

真実的なJN0-364試験問題集 &合格スムーズJN0-364トレーニング費用 | 素敵なJN0-364資格準備



安全で信頼できるウェブサイトとして、あなたの個人情報の隠しとお支払いの安全性を保障していますから、弊社のJuniperのJN0-364試験ソフトを安心にお買いください。弊社のPassTestは最大なるIT試験のための資料庫ですので、ほかの試験に興味があるなら、PassTestで探したり、弊社の係員に問い合わせたりすることができます。心よりご成功を祈ります。

誰もが知っているように、JuniperのJN0-364模擬テストシミュレーションは試験の成功に重要な役割を果たします。シミュレーションにより、JN0-364試験問題の無料デモを利用して、実際の試験の状況を把握できます。昔のことわざにあるように、敵とあなた自身を知っているので、敗北の危険なしに100回戦うことができます。PassTestのJN0-364トレーニング資料のシミュレーションにより、あなたの長所と短所を明確に理解できると同時に、JN0-364試験について包括的に学び、簡単にService Provider Routing and Switching, Specialist (JNCIS-SP)合格することができます。

>> JN0-364試験問題集 <<

ユニークJuniper JN0-364 | 更新するJN0-364試験問題集試験 | 試験の準備方法Service Provider Routing and Switching, Specialist (JNCIS-SP) トレーニング費用

Juniper完璧な素材の選択については、品質、価格、アフターサービスなどに反映される場合があります。JN0-364試験シミュレーションは、PassTest試験のシラバスに厳密に基づいた試験に関する知識の蓄積です。情報や試験へのアクセスをユーザーに提供し、教室のように参加したときに模擬的なテスト環境を提供します。また、JN0-364学習ガイドの内容は、日常生活での実践に適した専門家によって選択されます。JN0-364準備資料Service Provider Routing and Switching, Specialist (JNCIS-SP)を渡すのに十分な時間がない忙しい労働者にとって特に有利です。

Juniper Service Provider Routing and Switching, Specialist (JNCIS-SP) 認定 JN0-364 試験問題 (Q44-Q49):

質問 # 44

You are configuring LDP in a service provider network. After enabling LDP on core interfaces, you notice that labels are being advertised for every loopback IPv4 address that is in your IGP. Which label distribution mode is being used in this scenario?

- A. ordered control
- B. conservative retention
- C. downstream unsolicited
- D. downstream on demand

正解: C

解説:

In the context of the Label Distribution Protocol (LDP), the method by which a router advertises labels to its neighbors is defined by its Label Advertisement Mode. According to Juniper Networks documentation and industry standards (RFC 5036), there are two primary modes: Downstream Unsolicited (DU) and Downstream on Demand (DoD).

In Downstream Unsolicited (DU) mode, which is the default behavior for Junos OS and most service provider implementations, an LSR (Label Switching Router) does not wait for a specific request from its neighbors.

Instead, as soon as the LSR learns a prefix through its Interior Gateway Protocol (IGP) and establishes an LDP session, it automatically generates a label for that prefix and advertises it to all of its LDP peers. This explains the scenario where labels appear for every loopback address in the IGP as soon as LDP is enabled.

DU mode is highly efficient for fast convergence because the labels are already present in the neighbors' databases before they are even needed for traffic forwarding.

By contrast, Downstream on Demand (DoD) requires a router to explicitly request a label for a specific prefix from its next-hop neighbor. Ordered Control (Option B) and Independent Control refer to the timing of label creation (whether a router waits for the next-hop to provide a label before creating its own), while Conservative Retention (Option A) refers to how a router stores labels it receives but doesn't currently use for forwarding. In the Junos default environment, LDP utilizes Downstream Unsolicited advertisement combined with Ordered Control and Liberal Retention to ensure a robust and rapidly converging MPLS control plane.

質問 # 45

What prevents routing loops in a single-area OSPF network?

- A. Forwarding policies
- B. The Bellman-Ford algorithm
- C. The Dijkstra algorithm
- D. Routing policies

正解: C

解説:

In OSPF, loop prevention within a single area is achieved through the fundamental nature of its link-state architecture. Unlike distance-vector protocols that rely on "routing by rumor," OSPF ensures that every router within an area maintains an identical Link-State Database (LSDB). This database acts as a complete map of the network topology.

Once the LSDB is synchronized, each router independently executes the Shortest Path First (SPF) algorithm, which is formally known as the Dijkstra algorithm. This mathematical process treats the local router as the "root" of a tree and calculates the shortest path to every other node (router) and prefix in the area based on the cumulative interface costs. Because every router uses the same synchronized map (the LSDB) and the same deterministic algorithm, they all arrive at a consistent, loop-free view of the best paths.

According to Juniper Networks technical documentation, the Dijkstra algorithm is superior to the Bellman-Ford algorithm (used by distance-vector protocols like RIP) in this regard. Bellman-Ford is susceptible to

"count-to-infinity" problems and loops because routers only know the distance and direction to a destination provided by their neighbors, rather than the full topology. In OSPF, even if a link fails, the updated Link-State Advertisement (LSA) is flooded rapidly, and the Dijkstra algorithm is re-run to find a new loop-free path.

Routing policies (Option B) are used to manipulate path selection or filter routes but are not the primary mechanism for fundamental loop prevention in OSPF. Similarly, forwarding policies (Option D) govern how traffic is handled at the data plane level rather than determining the control plane's loop-free topology.

質問 # 46

Exhibit:

Referring to the exhibit, R1 and R2 are configured to run IS-IS. The IS-IS adjacency between R1 and R2 is up. What does the output of the show isis interface command tell you about R1?

- A. R1 advertises a Level 1 metric of 100 and a Level 2 metric of 100 toward R2 in its link-state PDU.
- B. R1 only forms a Level 2 adjacency with R2.
- C. R1 sends Level 1 hello PDUs to R2.
- D. R1 is not configured to use wide metrics.

正解: B

解説:

In the IS-IS (Intermediate System to Intermediate System) protocol as implemented in Junos OS, routers can operate at two

hierarchical levels: Level 1 (L1) for intra-area routing and Level 2 (L2) for inter-area backbone routing. By default, a Juniper router and its interfaces are configured to act as Level 1/2, meaning they will attempt to form adjacencies at both levels simultaneously.

According to Juniper Networks technical documentation, the `show isis interface` command provides a granular view of how the protocol is interacting with specific local links. In the provided exhibit, we must examine the L (Level) column and the DR (Designated Router) status columns to understand R1's operational state.

* Level Configuration: Under the L column for both the physical interface `ge-0/0/0.0` and the loopback `lo0.0`, the value is strictly 2. This indicates that these interfaces have been explicitly configured to operate only at Level 2.

* Adjacency Capabilities: For the interface `ge-0/0/0.0`, the Level 1 DR field is marked as Disabled. This confirms that R1 is not participating in Level 1 operations on this link; it will not transmit Level 1 Hello PDUs, nor will it listen for them. Consequently, R1 is incapable of forming a Level 1 adjacency with R2 on this segment.

* Metric Implications: The exhibit shows an L1/L2 Metric of 100/100. In Junos, "narrow" metrics (the default) are limited to a maximum value of 63 per interface. A metric of 100 indicates that wide metrics (wide-metrics-only) have been enabled. Therefore, option A is incorrect because the router is using wide metrics.

Since the prompt states the adjacency is "up," and the interface is restricted to Level 2, we can conclude that R1 only forms a Level 2 adjacency with R2 (Option B). Even though an L1 metric of 100 is displayed in the table as a configured value, it is not actually "advertised" in a Link-State PDU because the Level 1 protocol is disabled on that interface.

質問 # 47

Exhibit:

```
user@R2> show route 198.51.100.1
inet.0: 19 destinations, 19 routes (19 active, 0 holddown, 0 hidden)
Restart Complete
+ = Active Route, - = Last Active, * = Both
198.51.100.1/32 *[Static/5] 5d 21:02:26
> to 203.0.113.65 via ge-0/0/3.0
user@R2> show route 172.20.110.0/24
inet.0: 19 destinations, 19 routes (19 active, 0 holddown, 0 hidden)
Restart Complete
+ = Active Route, - = Last Active,
* = Both
172.20.110.0/24 *[Static/5] 10:43:01
> via gr-0/0/0.0
```

Referring to the exhibit, traffic destined to which network will be sent through the tunnel?

- A. 198.51.100.1/32
- B. 203.0.113.65
- C. 172.20.110.0/24
- D. 0.0.0.0/0

正解: C

解説:

To determine which traffic is being sent through a tunnel in a Junos OS environment, an administrator must analyze the routing table output for the exit interface associated with each destination prefix. The provided exhibit shows the results of the `show route` command on router R2 for two specific destination networks.

In the first output, the destination `198.51.100.1/32` is an active static route. The next-hop information specifies that traffic for this address is sent to the gateway `203.0.113.65` via the interface `ge-0/0/3.0`. According to Juniper Networks interface naming conventions, the prefix `ge-` denotes a Gigabit Ethernet interface, which represents a standard physical connection. Therefore, this traffic does not traverse a tunnel.

In the second output, the destination `172.20.110.0/24` is also an active static route. However, the next-hop for this network is listed as `via gr-0/0/0.0`. In the Junos operating system, the `gr-` prefix explicitly identifies a Generic Routing Encapsulation (GRE) tunnel interface. GRE is a widely used protocol in service provider networks to encapsulate various network layer protocols over an IP backbone, effectively creating a virtual point-to-point link. Because the routing table has installed the route for `172.20.110.0/24` specifically via the `gr-` interface, all traffic destined for this network will be encapsulated and sent through the tunnel.

The other choices are incorrect for the following reasons:

* `203.0.113.65` (Option B): This is the next-hop IP address for the physical Gigabit Ethernet path; it is not a destination network directed to a tunnel.

* `0.0.0.0/0` (Option C): There is no information in the exhibit regarding a default route.

* `198.51.100.1/32` (Option D): As identified by the `ge-` interface prefix in the exhibit, traffic for this destination is sent via a physical Ethernet link.

質問 # 48

Which IS-IS packet type will establish and maintain neighbor relationships?

- A. partial sequence number PDU
- B. update PDU
- C. link-state PDU
- D. hello PDU

正解: D

解説:

In the IS-IS (Intermediate System to Intermediate System) protocol, communication between routers is performed using Protocol Data Units (PDUs). To discover neighbors and maintain adjacencies, IS-IS relies on the Hello PDU (IIH - IS-IS Hello).

According to Juniper Networks technical documentation, when IS-IS is enabled on an interface, the router begins transmitting Hello PDUs to a multi-destination address (multicast). These PDUs contain essential information such as the router's System ID, its configured Area Addresses, and its Level capability (Level 1, Level 2, or both). For two routers to become neighbors, they must exchange these Hello PDUs and agree on specific parameters, such as the MTU of the link and the hello/hold timers.

Once an adjacency is established, the Hello PDU serves as a "keepalive" mechanism. If a router stops receiving Hello PDUs from a neighbor for a duration exceeding the Holding Time, it assumes the neighbor is down and flushes the associated Link-State PDUs (LSPs) from its database.

To clarify the other options:

* Link-State PDU (Option A): These are used to distribute actual topology and reachability information, not to form adjacencies.

* Partial Sequence Number PDU (Option C): PSNPs are used on point-to-point links to acknowledge the receipt of LSPs or to request missing LSPs.

* Update PDU (Option D): This is not a standard IS-IS term; in IS-IS, updates are handled via the flooding of LSPs.

質問 # 49

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最も短い時間で自分のIT技能を強化したいけれど、質の良い学習教材がないので悩んでいますか。ご心配なく、PassTestのJuniperのJN0-364試験トレーニング資料を手に入れるなら、ITに関する認定試験はなんでも楽に合格できます。PassTestのJuniperのJN0-364試験トレーニング資料は高度に認証されたIT領域の専門家の経験と創造を含めているものです。PassTestは君にとって、ベストな選択だといっても良いです。

JN0-364トレーニング費用: <https://www.passtest.jp/Juniper/JN0-364-shiken.html>

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京都きょうとに入はいって以来いらい、軍事ぐんじに市政しせいに日夜にちJN0-364や奔走ほんそうし、神経しんけいを休やすめ得えた日ひというのがない、見たところ椿を取り押さえている巨漢以外は雄介でも応戦可能だと判断できる。

実際の-最新のJN0-364試験問題集試験-試験の準備方法JN0-364トレーニング費用

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